COSC 240 Spring 2019 Problem Set 6 Due by class time on April 11, 2019 40 points

1. (20 points) Average consensus

Consider a set of n nodes (named 1, 2, ..., n) connected by an undirected graph G(V,E). Each node i maintains a local variable x_i , which is initialized to the input of node i. The nodes perform an iterative computation described here. The value of x_i after t iterations of the algorithm is denoted as x_i [t]. Thus, x_i [0] equals the input of node i.

In t-th iteration (or round) of the algorithm, each node obtains x_j [t-1] from each neighbor j, that is, from each node j such that edge (i,j) is in E.

In the following, assume that $a_{ij} = 0$ whenever edge (i,j) is not in E.

In the t-th iteration, each node i updates its x variable as follows.

$$x_i[t] = \sum_{i=1}^n a_{ji} x_j[t-1], \text{ for t>0}$$

Observe that in the above equation $a_{ij} = 0$ whenever edge (i,j) is not in E. Thus, to perform the above computation node i does NOT need x_j if j is not i's neighbor.

Now let us consider an n-by-n matrix A such that $A[i,j] = a_{ij}$

The above algorithm ensures that asymptotically the x values at all the nodes converge to the average of the inputs at all the nodes, provided that the following conditions are true:

- G(V,E) is a connected graph.
- $a_{ij} >= 0$ for all i,j
- The elements in each row of A add to 1.
- The elements in each column of A add to 1.
- $a_{ij} > 0$ if and only if (i,j) is in E.

Suppose that V={1,2,3} and E={(1,2), (2,3)}.

Suppose that the inputs at nodes 1, 2 and 3 are 4, 5, and 6, respectively.

(a) Will matrix A below result in average consensus? Explain why.

	7/8	1/8	ן 0
A =	1/8	1/4	5/8
	0	5/8	3/8]

- (b) Determine A² and A⁴.
- (c) Determine the values of variable x at the three nodes after 2 iterations and 4 iterations of the algorithm.Hint: You should be able to use the answer from part (b).
- (d) Should the x values at all the nodes add to 15 after t iterations, for any t? Explain why.
- (20 points) Consider the memoization algorithm for determining the length of the longest common subsequence, as given on slides 26-27 of <u>https://ocw.mit.edu/courses/electrical-engineering-and-computer-science/6-046jintroduction-to-algorithms-sma-5503-fall-2005/video-lectures/lecture-15-dynamicprogramming-longest-common-subsequence/lec15.pdf
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Using amortized analysis, argue that its execution time is O(mn). You may find it convenient to apply the accounting method to the recursion tree in this case, observing that a constant amount of work is performed (on average) at each node of the recursion tree.

Note: An informal argument will suffice.